Demand Response: Congestion in the Electricity Network

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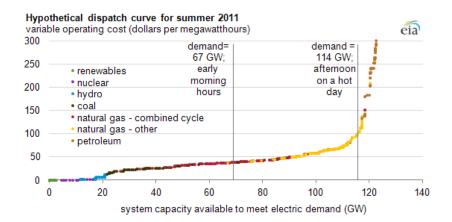
June 26, 2017

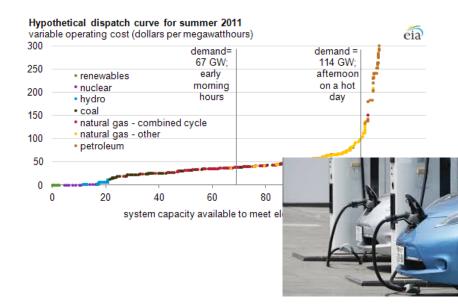
Summer School on Network Theory CIGNE, Roscoff

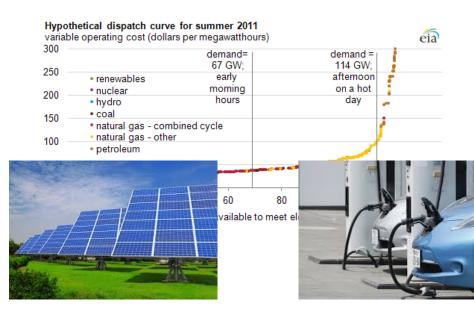
Paulin J. (EDF - Inria)

Demand Response: exploiting flexibilities

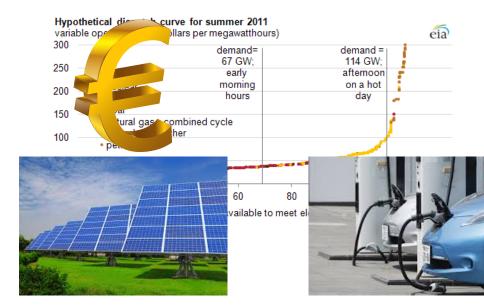
June 26, 2017 1 / 7







Introduction: Demand Response



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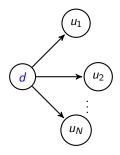


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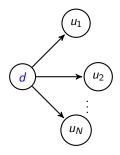
June 26, 2017 2 / 7

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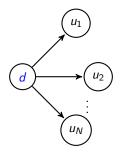
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- Consumers converge to an equilibrium consumption profile.

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$$\min_{\ell_n \in \mathbb{R}^T} (1 - \alpha) b_n(\ell_n, \ell_{-n}) - \alpha u_n(\ell) \quad (1a)$$
s.t.
$$\sum_{t \in \mathcal{T}} \ell_n^t = E_n, \quad (1b)$$

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Proposition

Assume
$$\forall n \in \mathcal{N}, \ \frac{\hat{\ell}_n^P}{E_n} + \frac{1}{2} \ge \frac{\hat{\ell}_p^P}{E}$$
, then, for $\alpha \in (0, 1]$, the NE of $\mathcal{G}_{\alpha}^{DP}$ gives:
 $\forall h \in \{P, O\}, \ \ell^h = E/2 + \alpha \times (\hat{\ell}^{\bar{h}} - \hat{\ell}^{\bar{h}})/2$. (2)

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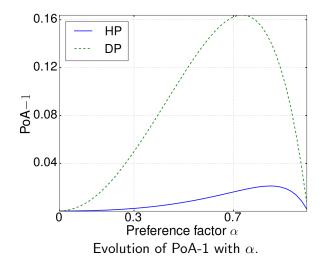
Assume
$$\forall n \in \mathcal{N}$$
, $\hat{\ell}_n^P \geq \frac{(\hat{\ell}^P - \hat{\ell}^O) - E_n}{2(N-1)}$, then $\forall \alpha \in [0, 1]$, the NE of $\mathcal{G}_{\alpha}^{HP}$ gives:

$$\forall h \in \{P, O\}, \ \ell^h = E/2 + \phi(\alpha) \times (\hat{\ell}^h - \hat{\ell}^{\bar{h}})/2 \ . \tag{3}$$

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where
$$\phi(\alpha) \stackrel{\text{def}}{=} \frac{2\alpha}{(1+\alpha)+(1-\alpha)N} \in [0,1].$$

Efficiency: Price of Anarchy



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THANK YOU!